

```

Y[i] = 2*X[i] mod 256
if (i>0 and X[i-1]>127) then Y[i]=Y[i]+1
end-for
if (X[15] > 127) then Y[0] = Y[0] xor 0x87

```

- 1 NOTE - Conceptually, the procedure is a left shift of each byte by one bit with carry propagating from one byte to the
2 next. Also, if the 15th (last) byte shift results in a carry, a special value (hexadecimal 0x87) is xor'ed into the first byte.
3 This value is derived from the modulus of the Galois Field (polynomial $x^{128}+x^7+x^2+x+1$).

4 5.2.2 EME2-AES Encryption

5 The EME2-AES encryption procedure can be described by the formula:

$$6 \quad C = \text{EME2-AES-Enc}(Key, T, P),$$

7 where

8 *Key* is the 48 or 64 byte EME2-AES key

9 *T* is the value of the associated data, of arbitrary byte length (zero or more bytes)

10 *P* is the plaintext, of length 16 bytes or more

11 *C* is the ciphertext resulting from the operation, of the same byte-length as *P*

12 The input to the EME2-AES encryption procedure is parsed as follows:

- 13 — The key is partitioned into three fields, $Key = \underline{KEY_{AD}} \mid \underline{KEY_{ECB}} \mid \underline{KEY_{AES}}$, with KEY_{AD} (the
14 associated data key) consisting of the first 16 bytes, KEY_{ECB} (the ECB pass key) consisting of the
15 following 16 bytes, and KEY_{AES} (the AES encryption/decryption key) consisting of the remaining
16 16 or 32 bytes.
- 17 — If not empty, the associated data *T* is partitioned into a sequence of blocks $T = T_1 \mid T_2 \mid \dots \mid T_r$,
18 where each of the blocks T_1, T_2, \dots, T_{r-1} is of length exactly 16 bytes, and T_r is of length between 1
19 and 16 bytes.
- 20 — The plaintext *P* is partitioned into a sequence of blocks $P = P_1 \mid P_2 \mid \dots \mid P_m$, where each of the
21 blocks P_1, P_2, \dots, P_{m-1} is of length exactly 16 bytes, and P_m is of length between 1 and 16 bytes.
- 22 The ciphertext shall then be computed by the sequence of steps in [Table 2 and Table 2](#) or equivalent. An
23 illustration of these steps (for plaintext of 130 full blocks and one partial block) is provided in [Figure 2](#).

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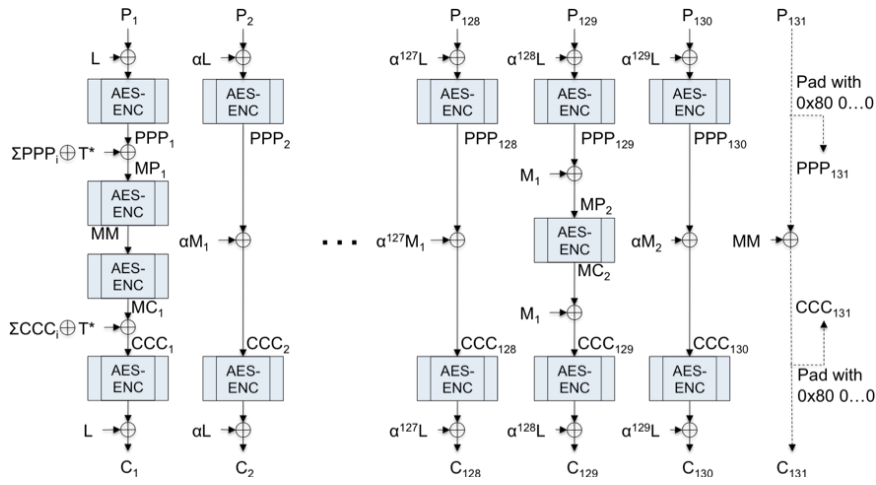


Figure 2, An illustration of EME2-AES: T^* is computed from the associated data using the function H from Table 2, and the M_i 's are computed as $M_i = MP_i \text{ xor } MC_i$

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Table 1, The function H for processing the associated data T

```

// Function H( KEYAES, KEYAD, T = (T1 ... Tr-1 Tr) ) :
1. if len(T) == 0 then
2.   T star = AES-Enc(KEYAES, KEYAD)
3. else
4.   KEYAD = Mult-by-alpha(KEYAD)
5.   for j = 1 to r - 1
6.     TTj = AES-Enc(KEYAES, KEYAD ⊕ Tj) ⊕ KEYAD
7.     KEYAD = Mult-by-alpha(KEYAD)
8.     if len(Tj) < 16 then
9.       Tj = Tj | 0x80 0x00... // pad Tj to 16 bytes, 0x80 followed by 0's
10.    KEYAD = Mult-by-alpha(KEYAD)
11.    TTr = AES-Enc(KEYAES, KEYAD ⊕ Tr) ⊕ KEYAD
12.    T star = TT1 ⊕ TT2 ⊕ ... ⊕ TTr

13. return T star // return the 16 byte value T star
    
```

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Table 2, The EME2-AES Encryption Procedure

```

// Function EME2-AES-Enc( KEYAES, KEYECB, KEYAD, T, P = (P1 ... Pm-1 Pm) ) :
1. T star = H( KEYAES, KEYAD, T) // Process the associated data
2. if len(Pm) = 16 then
3.   lastFull = m
4. else
5.   lastFull = m - 1
6.   PPPm = Pm | 0x80 0x00 ... // Pad Pm to 16 bytes, 0x80 followed by 0's

// First ECB pass
7. L = KEYECB
8. for i = 1 to lastFull
9.   PPPi = AES-Enc( KEYAES, L ⊕ Pi )
10.  L = Mult-by-alpha(L)

// Intermediate mixing
11. MP = PPP1 ⊕ PPP2 ⊕ ... ⊕ PPPm ⊕ T star
12. if len(Pm) < 16 then
13.   MM = AES-Enc( KEYAES, MP )
14.   MC = MC1 = AES-Enc( KEYAES, MM )
15. else
16.   MC = MC1 = AES-Enc( KEYAES, MP )
17. M = M1 = MP ⊕ MC
18. for j = 2 to lastFull
19.   if (j - 1 mod 128 > 0) then // use the current mask M
20.     M = Mult-by-alpha(M)
21.     CCCj = PPPj ⊕ M
22.   else // calculate a new mask M
23.     MP = PPPj ⊕ M1
24.     MC = AES-Enc( KEYAES, MP )
25.     M = MP ⊕ MC
26.     CCCj = MC ⊕ M1
27. if lastFull < m then
28.   Cm = Pm ⊕ (MM truncated to len(Pm) bytes)
29.   CCCm = Cm | 0x80 0x00 ... // Pad Cm to 16 bytes, 0x80 followed by 0's
30. CCC1 = MC1 ⊕ CCC2 ⊕ ... ⊕ CCCm ⊕ T star
    
```

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```

// Second ECB Pass
31. L = KEY_ECB
32. for i = 1 to lastFull
33.   Ci = AES-Enc(KEY_AES, CC) ⊕ L
34.   L = Mult-by-alpha(L)
35. return C = (C1 ... Cm-1 Cm)
    
```

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2 **5.2.3 EME2-AES Decryption**

3 The EME2-AES decryption procedure can be described by the formula:

4 $C = \text{EME2-AES-Dec}(Key, T, C)$,

5 where

6 *Key* is the 48 or 64 byte EME2-AES key

7 *T* is the value of the associated data, of arbitrary byte length (zero or more bytes)

8 *C* is the ciphertext, of length 16 bytes or more

9 *P* is the plaintext resulting from the operation, of the same byte-length as *C*

10 The input to the EME2-AES decryption procedure is parsed as follows:

- 11 — The key is partitioned into three fields, $Key = KEY_{AD} | KEY_{ECB} | KEY_{AES}$, with KEY_{AD} (the associated data key) consisting of the first 16 bytes, KEY_{ECB} (the ECB pass key) consisting of the following 16 bytes, and KEY_{AES} (the AES encryption/decryption key) consisting of the remaining 16 or 32 bytes.
- 12 — If not empty, the associated data is partitioned into a sequence of blocks $T = T_1 | T_2 | \dots | T_r$, where each of the blocks T_1, T_2, \dots, T_{r-1} is of length exactly 16 bytes, and T_r is of length between 1 and 16 bytes.
- 13 — The ciphertext *P* is partitioned into a sequence of blocks $C = C_1 | C_2 | \dots | C_m$, where each of the blocks C_1, C_2, \dots, C_{m-1} is of length exactly 16 bytes, and C_m is of length between 1 and 16 bytes.

14 The plaintext shall then be computed by the sequence of steps in Table 3 or equivalent. (Note that the only difference between the encryption and decryption procedures is that all the AES-Enc calls in Table 3 are replaced in Table 4 by calls to AES-Dec. Also note that the function H is defined in Table 1).

15 Table 3, The EME2-AES Decryption procedure

```

// Function EME2-AES-Dec(KEY_AES, KEY_ECB, KEY_AD, T, C = (C1 ... Cm-1 Cm)) :
    
```

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```

1.  $T_{star} = H(KEY_{AES}, KEY_{AD}, T)$  // Process the associated data
2. if  $len(C_m) = 16$  then
3.    $lastFull = m$ 
4. else
5.    $lastFull = m - 1$ 
6.    $CCC_m = C_m | 0x80 0x00 \dots$  // Pad  $C_m$  to 16 bytes, 0x80 followed by 0's

// First ECB pass
7.  $L = KEY_{ECB}$ 
8. for  $j = 1$  to  $lastFull$ 
9.    $CCC_j = AES-Enc(KEY_{AES}, L \oplus C_j)$ 
10.   $L = \text{Mult-by-alpha}(L)$ 

// Intermediate mixing
11.  $MC = CCC_1 \oplus CCC_2 \oplus \dots \oplus CCC_m \oplus T_{star}$ 
12. if  $len(C_m) < 16$  then
13.    $MM = AES-Dec(KEY_{AES}, MC)$ 
14.    $MP = MP_1 = AES-Dec(KEY_{AES}, MM)$ 
15. else
16.    $MP = MP_1 = AES-Dec(KEY_{AES}, MC)$ 
17.    $M = M_1 = MP \oplus MC$ 
18.   for  $j = 2$  to  $lastFull$ 
19.     if  $(j - 1 \bmod 128 > 0)$  then // use the current mask M
20.        $M = \text{Mult-by-alpha}(M)$ 
21.        $PPP_j = CCC_j \oplus M$ 
22.     else // calculate a new mask M
23.        $MC = CCC_j \oplus M_1$ 
24.        $MP = AES-Dec(KEY_{AES}, MC)$ 
25.        $M = MP \oplus MC$ 
26.        $PPP_j = MP \oplus M_1$ 
27.   if  $lastFull < m$  then
28.      $P_m = C_m \oplus (MM \text{ truncated to } len(C_m) \text{ bytes})$ 
29.      $PPP_m = P_m | 0x80 0x00 \dots$  // Pad  $P_m$  to 16 bytes, 0x80 followed by 0's
30.      $PPP_1 = MP_1 \oplus PPP_2 \oplus \dots \oplus PPP_m \oplus T_{star}$ 

// Second ECB Pass
31.  $L = KEY_{ECB}$ 
32. for  $j = 1$  to  $lastFull$ 
33.    $P_j = AES-Dec(KEY_{AES}, PPP_j \oplus L)$ 
34.    $L = \text{Mult-by-alpha}(L)$ 

35. return  $P = (P_1 \dots P_{m-1} P_m)$ 

```

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